

Some Optimization Problems in Quantum Information Science

Chi-Kwong LI

(Ferguson Professor) College of William and Mary, Virginia,
(Affiliate member) Institute for Quantum Computing, Waterloo

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- In quantum science, one needs to manipulate **quantum states** using **quantum operations**.
- One may also want to estimate the change of a quantum states after they go through a certain **quantum channel**.

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Example Suppose $\mathcal{F}_1, \mathcal{F}_2 \subseteq M_2$. Then ...

Some Known Results

- (Chefles, Jozsa, Winter, 2004) $\mathcal{F}_1, \mathcal{F}_2$ are families of pure states

$$\rho_i = x_i x_i^* \text{ and } \sigma_i = y_i y_i^* \text{ for } i = 1, \dots, k.$$

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- (Huang, Li, E. Poon, Sze, 2012) General families. Solve some complicated matrix equations.

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$$F(\rho_1, \rho_2) = \|\sqrt{\rho_1}\sqrt{\rho_2}\| \leq \|\sqrt{\sigma_1}\sqrt{\sigma_2}\|. \quad (1)$$

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$$\sigma = \frac{1}{|\mu_1\mu_2|}(\sigma_3 - |\mu_1|^2\sigma_1 - |\mu_2|^2\sigma_2) = \operatorname{Re}\sqrt{\sigma_1}C\sqrt{\sigma_2}$$

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Question Can we find a more explicit (and symmetric) conditions on x_1, x_2, x_3 , and $\sigma_1, \sigma_2, \sigma_3$ for the existence of Φ ?

- (Choi, 1975) A linear operator $\Phi : M_n \rightarrow M_m$ is a quantum operation if and only if the (Choi) matrix $P = (\Phi(E_{ij}))_{1 \leq i, j \leq n} \in M_n(M_m)$ is positive semi-definite with $\text{tr} \Phi(E_{ij}) = \delta_{ij}$.

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- (D. Drusvyatskiy, C.K. Li, D. Pelejo, Y.L. Voronin, H. Wolkowicz, 2015) General families. Construct a Choi matrix $P = (P_{ij}) \in M_n(M_m)$ such that

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- Evidently,

$$\begin{aligned} \{\text{Unitary operations}\} &\subseteq \{\text{Mixed unitary operations}\} \\ &\subseteq \{\text{Unital operations}\} \subseteq \{\text{General quantum operations}\}. \end{aligned}$$

Approximation problems

Question For a given $\varepsilon > 0$, determine whether there is a quantum operation Φ such that $\|\Phi(\rho_j) - \sigma_j\| < \varepsilon$ for all $j = 1, \dots, k$.

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the **Frobenius norm** $\|X\|_{\text{Fr}} = (\text{tr}(X^*X))^{1/2}$.

Approximation

- There are results on the upper bound and lower bounds for $d(\Phi(\rho_1), \sigma_1)$ for $\Phi \in \mathcal{S}$, where

\mathcal{S} is the set of all unitary, mixed unitary, unital, or general channels, and

$d(\alpha, \beta)$ are different measures such as

$\|\alpha - \beta\|$ for a unitary similarity invariant norm $\|\cdot\|$,

the Fidelity function $d(\alpha, \beta) = F(\alpha, \beta)$,

the relative entropy function $d(\alpha, \beta) = S(\alpha||\beta)$.

- We first describe results on $\mathcal{F}_1 = \{A\}$ and $\mathcal{F}_2 = \{B\}$.

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Theorem

Let $\|\cdot\|$ be a USI norm, σ_1, ρ_1 are density matrices with eigenvalues

$$a_1 \geq \cdots \geq a_n \quad \text{and} \quad b_1 \geq \cdots \geq b_n.$$

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Let $\rho, \sigma \in M_n$ be density matrices. The following are equivalent.

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④ $\lambda(\sigma) \prec \lambda(\rho)$, i.e., the sum of the k largest eigenvalues of σ is not larger than that of ρ for $k = 1, \dots, n - 1$.

Theorem (based on a result in [Li & Tsing, 1989])

Let $\|\cdot\|$ be a USI norm, σ_1, ρ_1 are density matrices with eigenvalues

$$a_1 \geq \cdots \geq a_n \text{ and } b_1 \geq \cdots \geq b_n.$$

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Step 2. Let $2 \leq j < k \leq \ell \leq n$ be such that

$$\Delta_{j-1} \neq \Delta_j = \cdots = \Delta_{k-1} < \Delta_k = \cdots = \Delta_\ell \neq \Delta_{\ell+1}.$$

Replace each $\Delta_j, \dots, \Delta_\ell$ by $(\Delta_j + \cdots + \Delta_\ell)/(\ell - j + 1)$, and go to Step 1.

Examples

Here are two examples illustrating the algorithm in the theorem.

Example 1 Let $\sigma_1 = \frac{1}{10}\text{diag}(4, 3, 3, 0)$ and $\rho_1 = \frac{1}{10}\text{diag}(3, 3, 3, 1)$.

Apply Step 0:

$$\text{Set } (\Delta_1, \dots, \Delta_4) = \frac{1}{10}\text{diag}(4, 3, 3, 0) - \frac{1}{10}\text{diag}(3, 3, 3, 1) = \frac{1}{10}\text{diag}(1, 0, 0, -1).$$

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where d_1, \dots, d_n are determined as follows.

Step 0. Suppose $a_1 \geq \dots \geq a_r \geq 0 = a_{r+1} = \dots = a_n$. Let

$$a = (a_1, \dots, a_r), \quad b = (b_1, \dots, b_r), \quad (d_{r+1}, \dots, d_n) = (b_{r+1}, \dots, b_n).$$

Go to Step 1.

Step 1. Let $k \in \{1, \dots, r\}$ be the largest positive integer such that

$$\frac{1}{a_1 + \dots + a_k} (a_1, \dots, a_k) \prec \frac{1}{b_1 + \dots + b_k} (b_1, \dots, b_k).$$

Set

$$(d_1, \dots, d_k) = \frac{a_1 + \dots + a_k}{b_1 + \dots + b_k} (a_1, \dots, a_k).$$

If $k = r$, then exit. Else, replace r, a, b by $r - k, (a_{k+1}, \dots, a_r), (b_{k+1}, \dots, b_r)$ and go to Step 1.

Examples If $(a_1, \dots, a_n) \prec (b_1, \dots, b_n)$, then $(d_1, \dots, d_n) = (a_1, \dots, a_n)$.

If $(b_1, \dots, b_n) = (1/n, \dots, 1/n)$, then $(d_1, \dots, d_n) = (1/n, \dots, 1/n)$.

More results and questions

- We also obtained results for general quantum channels, and other functions on two density matrices such as the **relative entropy**:

$$S(\rho_1 || \rho_2) = \text{tr } \rho_1 (\log_2 \rho_1 - \log_2 \rho_2).$$

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- Minimize/maximize $d((\Phi(\rho_1), \dots, \Phi(\rho_k)), (\sigma_1, \dots, \sigma_k))$ for other distance measure d ?
- one may start with the study of $\|\Phi(\rho_1 + i\rho_2) - (\sigma_1 + i\sigma_2)\|$ for the a special norm.

Thank you for your attention!